

- For any regular language, there is a unique *minimal* DFA
- By finding the minimal DFA, we can also prove equivalence (or not) of different FSA and the languages they recognize
- In general the idea is:
 - Throw away unreachable states (easy)
 - Merge equivalent states
- There are two well-known algorithms for minimization:
 - Hopcroft's algorithm: find and eliminate equivalent states by partitioning the set of states
 - Brzozowski's algorithm: 'double reversal'

Finding equivalent states

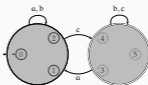
Intuition



The edges leaving the group of nodes are identical.
Their *right languages* are the same.

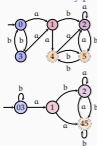
Finding equivalent states

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Minimization by partitioning



- Accepting & non-accepting states form a partition
 $Q_1 = \{0, 1, 2, 3\}, Q_2 = \{4, 5\}$
- For any of the nodes in a set, if any symbol leads to different sets of nodes, split
- $Q_1 = \{0, 3\}, Q_2 = \{1\}, Q_3 = \{2, 4, 5\}$
- Stop when we cannot split any of the sets, merge the indistinguishable states

Minimization by partitioning

tabular version



- Create a state-by-state table, mark distinguishable pairs: (q_1, q_2) such that $\{ \Delta(q_1, x), \Delta(q_2, x) \}$ is a distinguishable pair for any $x \in \Sigma$



Minimization by partitioning

tabular version



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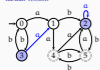


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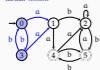


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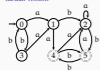
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- Merge indistinguishable states

Minimization by partitioning

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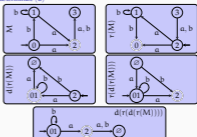
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- Merge indistinguishable states
- The algorithm can be improved by choosing which cell to visit carefully

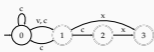
Brzozowski's algorithm

double reverse (r), determinize (d)



An exercise

find the minimum DFA for the automaton below



Minimization algorithms

final remarks

- There are many versions of the 'partitioning' algorithm. General idea is to form equivalence classes based on right-language of each state.
- Partitioning algorithm has $O(n \log n)$ complexity
- 'Double reversal' algorithm has exponential worst-time complexity
- Double reversal algorithm can also be used with NFAs (resulting in the minimal equivalent DFA - NFA minimization is intractable)
- In practice, there is no clear winner, different algorithms runs faster on different input
- Reading suggestion: [hopcroft1979](#), [jurfafsky2009](#)

Next:

- PST
- PSA and regular languages

Acknowledgments, credits, references

